Southern Illinois University Carbondale [OpenSIUC](http://opensiuc.lib.siu.edu?utm_source=opensiuc.lib.siu.edu%2Fece_confs%2F82&utm_medium=PDF&utm_campaign=PDFCoverPages)

[Conference Proceedings](http://opensiuc.lib.siu.edu/ece_confs?utm_source=opensiuc.lib.siu.edu%2Fece_confs%2F82&utm_medium=PDF&utm_campaign=PDFCoverPages)

[Department of Electrical and Computer](http://opensiuc.lib.siu.edu/ece?utm_source=opensiuc.lib.siu.edu%2Fece_confs%2F82&utm_medium=PDF&utm_campaign=PDFCoverPages) [Engineering](http://opensiuc.lib.siu.edu/ece?utm_source=opensiuc.lib.siu.edu%2Fece_confs%2F82&utm_medium=PDF&utm_campaign=PDFCoverPages)

3-2005

Distributed Detection with Channel Errors

M. Madishetty *Southern Illinois University Carbondale*

V. Kanchumarthy *Southern Illinois University Carbondale*

R. Viswanathan *Southern Illinois University Carbondale*, viswa@engr.siu.edu

Chandrakanth H. Gowda

Tuskegee University Follow this and additional works at: [http://opensiuc.lib.siu.edu/ece_confs](http://opensiuc.lib.siu.edu/ece_confs?utm_source=opensiuc.lib.siu.edu%2Fece_confs%2F82&utm_medium=PDF&utm_campaign=PDFCoverPages)

Published in Madishetty, M., Kanchumarthy, V., Viswanathan, R., & Gowda, C.H. (2005). Distributed detection with channel errors. Proceedings of the Thirty-Seventh Southeastern Symposium on System Theory, 2005 (SSST '05), 302-306. doi: 10.1109/SSST.2005.1460926 ©2005 IEEE. Personal use of this material is permitted. However, permission to reprint/republish this material for advertising or promotional purposes or for creating new collective works for resale or redistribution to servers or lists, or to reuse any copyrighted component of this work in other works must be obtained from the IEEE. This material is presented to ensure timely dissemination of scholarly and technical work. Copyright and all rights therein are retained by authors or by other copyright holders. All persons copying this information are expected to adhere to the terms and constraints invoked by each author's copyright. In most cases, these works may not be reposted without the explicit permission of the copyright holder.

Recommended Citation

Madishetty, M.; Kanchumarthy, V.; Viswanathan, R.; and Gowda, Chandrakanth H., "Distributed Detection with Channel Errors" (2005). *Conference Proceedings.* Paper 82. [http://opensiuc.lib.siu.edu/ece_confs/82](http://opensiuc.lib.siu.edu/ece_confs/82?utm_source=opensiuc.lib.siu.edu%2Fece_confs%2F82&utm_medium=PDF&utm_campaign=PDFCoverPages)

This Article is brought to you for free and open access by the Department of Electrical and Computer Engineering at OpenSIUC. It has been accepted for inclusion in Conference Proceedings by an authorized administrator of OpenSIUC. For more information, please contact [opensiuc@lib.siu.edu.](mailto:opensiuc@lib.siu.edu)

Distributed Detection with Channel Errors

M. **Madishetty, V.** Kanchumarthy, R. **Viswanathan** Department of Electrical & Computer Engineering Southem Illinois University Carbondale Carbondale IL 62901-6603

Chandrakanth H. **Gowda** Department of Electrical Engineering Tuskegee University Tuskegee, *AL,* **36088.**

Absfract-Performances of distributed detection (DD) systems employing a set of geographically dispensed sensors have been investigated for the past two decades. In this paper we study the variations in the false alarm **and detection probabilities of a DD system due to the errors caused by the links between sensors and the fusion center. Both asymptotic and finite sample performances are studied. The results bring out the exact dependence of these probabilities on the link reliability. Such a study is meaningfd because of the recent research interests in wireless sensor networks.**

I. **INTRODUCTION**

In **a** wireless sensor network with distributed **sensors,** each sensor makes measurement **with** regard to a phenomenon of interest (POI) in order to make a decision **on** the presence or the absence of **POI.** The **POI** might be a biological spill or **the** sighting of a vehicle of **an** adversary. Each sensor processes its own information and passes the condensed information to a cluster head (or fusion center) through a wireless channel. The **data** arriving **from** various sensors are **fused** together appropriately in order to **make** a **fmal** decision on the presence or the absence of a **POI.**

Several papers have addressed a myriad of signal processing issues in sensor networks **[la].** A number *of* problems **in** detection, classification, and **tracking** of targets are discussed in *[I].* The allocation of optimum number of quantization bits at the sensors for **a** rate constrained communication channel and a large number of sensors was addressed in **(21.** References **[34]** discuss **the** performances of different fusion rules that could be formulated based on varied knowledge of communication channel statistics. Another asymptotic (large number *of* **sensors)** optimization of wireless sensor networks for decentralized detection was addressed in *[5].* A relatively old contribution **[6-71** on optimal detection with faulty processors has relevance to distributed detection problem **with** channel errors. Binary symmetric channel was considered as the model for the sensor-to-fusion center link in **[8],** but the emphasis was on the person-by-person optimization of local sensor rule and the fusion rule.

Due to the hostile nature of **a** wireless channel, a sensor **data** might not be received reliably at the fusion center.

Hence, sensor quantization rules designed *for* a specific false alarm probability would not produce a fixed false alarm probability at the fusion center. **hi** this study we derive analytical expressions for the false alarm and detection performances of a distributed detection system **at** the fusion center **and examine** how much variation of false alarm probability can be anticipated. **To** our knowledge, a **study** of the variation of false **alarm** and detection probabilities due **to** changes in channel statistics has not been addressed in the literature. A study of the changes **in** these probabilities due to **a** randomized **data** selection strategy **was** addressed in **[9].**

11. EFFECT OF CHANNEL ERRORS ON THE RELIABILITY OF SENSOR **DECISION**

Consider **a** wireless sensor network consisting of *N* sensors. The **network** is deployed to assess the presence or the absence of **a** phenomenon of interest **(POI)** in a geographical area of interest. Sensor *i* gathers information **pertaining to the POI and makes a decision** u_i **(** $u_i = 1$ **for** deciding the presence of POI and $u_i = 0$ otherwise). Each sensor sends its decision **to a** fusion **center through** a communication link, which is not totally reliable. Let u_{i0} denote the decision of the *i*th sensor as received at the fusion center. **Let**

$$
P_{f,i} = P\left(u_i = 1 | \text{POI absent }\right), P_{si} = P\left(u_i = 1 | \text{POI present }\right),
$$

$$
P_{f,i}^c = P\left(u_{i0} = 1 | \text{POI absent }\right), P_{si}^c = P\left(u_{i0} = 1 | \text{POI present }\right),
$$

 $P_{c,i}$ = Probability of bit error of the i^{th} link. Assuming that the link performance is statistically independent of **the** decision made by the sensor, the reliability parameters of the sensor decision as received **by** the fusion center can be describe by the following set of equations:

$$
P_{f_i}^c = P_{f_i} (1 - P_{c_i}) + (1 - P_{f_i}) P_{c_i}
$$
 (1)

$$
P_{a1}^{c} = P_{a1}(1 - P_{c1}) + (1 - P_{a1})P_{c1}
$$
 (2)

Rewriting the above **equations** yields the following results:

$$
P_{f}^{c} = P_{f} (1 - 2 P_{c}) + P_{c} \tag{3}
$$

$$
P_{a}^{c} = P_{ai} (1 - 2 P_{ci}) + P_{ci}
$$
 (4)

Therefore, the reliability of the decision received u_{i0} could be different from that of the decision u_i made at the sensor.

0-7803-8808-9/05/\$20.00 *82005* **IEEE 302**

.

Assume that the link bit error, $P_{c,i}$ < 1/2 (if it is greater than $\frac{1}{2}$, then the decision rule of receiver for the i^{th} link at the fusion center could be complimented to achieve it to be less than $\frac{1}{2}$. If $P_{f,i} \leq 1/2$ then $P_{f,i}^c > P_{f,i}$. That is, the false alarm probability of the decision received at the fusion center is **higher** than the false alarm probability of the decision made **by** the sensor. *As* the link becomes very unreliable, both the link bit error probability and the probability, p^e , approach

%. **Similirly,** when the probability of detection at the sensor is greater than $\frac{1}{2}$, the detection probability of the decision received at the fusion center is less than the detection probability at the sensor. Only when $P_d \leq 1/2$, the link error "increases" the probability of detection, p_{a}^{c} , to be above that of $P_{d,i}$ (of course this is achieved with a concomitant increase in the false alarm probability). Given the unreliable nature of the communication link between a sensor and the fusion center, we examine in the next section its impact on the reliability of **the** decision made by the fusion center.

nI. PERFORMANCE **OF** THE **FUSION** CENTER

Let us assume that each sensor in the network makes a decision independent of others such that each exhibits **an** identical performance. That is, each sensor decision is independent and identically distributed given the true state of nature with regard to the presence or the absence of **POI. Also,** assume that each link between a sensor and the fusion center exhibits, **on an** average, an identical link error performance. Given **a** large **number** of sensors in the network, it is well **known** that, under very general conditions, **an** optimum fusion rule for combining the decisions **from** the sensors takes the form of a **counting** rule. That is, the fusion center declares that POI is present when the number of sensors declaring that the **POI** is present exceeds a certain pre-determined threshold *t*. We first examine what choice of *t* would be reasonable for the asymptotic condition of large ($N \rightarrow \infty$) number of sensors. Because of the assumptions mentioned above, hereafter we can drop the subscript *i* that identifies a particular sensor.

A. Asymptotic Condition:

When the number of sensors is large, we can apply the Gaussian approximation to the sum of binomial probabilities. Let P_{F0} *,P_{D0}* denote the false alarm probability and the detection probability, respectively, of the final decision made by the fusion center **and** let the count anived at the fusion center be denoted by $Z = \sum_{i=0}^{N} u_i$. Since the false alarm probability of the decision received **from** a sensor at the fusion center depends on the link error probability, it is reasonable **to** assume that its value can be bounded below an upper bound corresponding to **a** minimum reliability of the communication link. If we do not impose such a minimum **i=l**

reliability measure, then the false alarm probability **couid**

approach $\frac{1}{2}$, in the worst scenario, as indicated in the previous section. Denoting $P'_i = \alpha < \alpha$, we can write

$$
P_{F0} = P(Z \ge t \, | \text{POI absent}) = Q\left(\frac{t - N\alpha}{\sqrt{N\alpha(1-\alpha)}}\right) (5)
$$

If we let

$$
t = N \alpha_u, \tag{6}
$$

$$
P_{r_0} = Q\left(\sqrt{N} \frac{\alpha_s - \alpha}{\sqrt{\alpha(1-\alpha)}}\right) \tag{7}
$$

tends to zero as $N \to \infty$, as long as $\alpha_n > \alpha > 0$. Moreover,

$$
P_{\nu 0} = Q \left(-\sqrt{N} \frac{\beta - \alpha_s}{\sqrt{\beta(1 - \beta)}} \right) \tag{8}
$$

tends to 1 as $N \to \infty$, provided $\beta = p_r^c > \alpha_s$. If $\beta < \alpha_u$, then P_{D0} tends towards zero, and hence it is important that α be satisfied. This can be guaranteed as long as the signal-to-noise ratio *(SNR)* at the sensor is above a certain minimum value and the link bit emor rate is below a certain value. For example, when detecting a constant signal in AWGN, the detection probability **and** the false alarm probability **at** a sensor **are** related by

$$
P_d = Q(Q^{-1}(P_f) - \sqrt{SNR})
$$
\n(9)

Using (3) , (4) and (9) the required values of SNR and P_c to guarantee $\beta > \alpha_{\kappa}$ can be arrived at.

B. Finite N:

If *N* is only finite, then the above asymptotic results are not valid. Moreover, the false alarm probability at the fusion, for a designed value of *t*, could increase to a large unacceptable value as the sensor-to-fusion link becomes unreliable. We next examine the variations in the fusion rule performance as a function of the reliability of the sensor-to-fusion link. Towards this effort we characterize the link to be the result of **transmitting** orthogonal **binary FSK** signals in **a Rayleigh** fading channel. For a wireless sensor network, the Rayleigh fading channel is **an** appropriate model to assume. For simplicity, we do not **assume** any error control coding for this link. Certainly, an error control code would make the link more reliabie. But, to **remind** the reader, **our** aim here is to observe the impact of a less reliable link on the performance of the fusion rule. Using standard results for noncoherent detection of binary **FSK** in slow Rayleigh fading channels [IO], we **can** write the following relations:

$$
P_c = \frac{1}{2 + \gamma_o} \tag{10}
$$

$$
\alpha = P_f + \frac{1 - 2P_f}{2 + \gamma_o} \tag{11}
$$

303

$$
\beta = P_d + \frac{1 - 2P_d}{2 + \gamma_0} \tag{12}
$$

where γ_0 is the average SNR of the Rayleigh fading channel.

Above, we used an average value of P_o averaged with respect to the fading distribution. We show in **111.** *C* that, for independent trials and a counting rule at the fusion center, this is a correct procedure for calculating the overall fusion probability of False Alarm (Detection).

In Figures 1-4 we show the variations of P_{F0} and P_{D0} as a function of γ_0 for some values of P_f P_d *t* and *N*. The γ_0 **axis** (average channel SNR) shows values fiom a In Figures 1-4 we show the variations of P_{F0} and P_{D0} as a function of γ_0 for some values of P_L P_d *t* and *N*. The γ_0 axis (average channel SNR) shows values from a minimum value, which guarantees that γ_0 approaches infinity, α approaches P_f and P_{F0} approaches the value **that** would be obtained had the **links been** error he. Depending on the values of P_f , *t*, and *N*, the fusion false alarm probability could be two or three decades higher **than the** desired value. This is in contrast to the asymptotic case where perfect detection ($p_{r0} \rightarrow 0$, $p_{p0} \rightarrow 1$) is possible. For finite *N,* it is essential that the link reliability is greater **than** a certain minimum value in order that an acceptable P_{ε_0} is achieved. In general, except for weak sensor signal **conditions,** the effect of link errors on **the** detection probability is less severe, because the **sensor** detection probability will be larger than **0.1.** Notice that $P(Z \geq t | \text{POI present})$ is a monotonic increasing function of β . Hence, interestingly, when P_d <0.5, better detection probability, $P_{p,q}$, is achieved when the link is less reliable!

C. False Alarm Probability of Fused Decision With Independent but Identical Fading Links

(see **(4)).**

In this section we provide a proof to show that the average link error probability can be used for each link while computing the overall false alarm (and detection) error probability. Let $\gamma_1, \gamma_2, ..., \gamma_N$ be the instantaneous SNR of the received signal corresponding to the individual Iinks between a sensor and the fusion center and let P_f be the false alarm probability of the decision made **by** a sensor. Then, for a specified counting rule at the fusion center, the false alarm probability of the fused decision **can** be **written** as $P_{F0} = E[f(P_f, \gamma_1, \gamma_2, ..., \gamma_N, t] \gamma_1, \gamma_2, ..., \gamma_N)$ (13) where the expectation operation with respect to the distribution of the instantaneous SNRs and $f(.)$ describes the function that determines the conditional false alarm probability **for** a given **fusion** rule, conditioned on the **instantaneous SNRs.** For a counting rule with threshold **f,**

$$
P_{F0} = E\left(\sum_{i=1}^{N} P\left(\text{successes out of } N \text{ trials with } p_i\right) \right) \qquad (14)
$$

as the success probability for the j^{th} trial

where p_j depends on γ_j and P_f . Rewriting (14) yields

$$
P_{F0} = E\left(\sum_{\substack{s=t_{j_1+1}}^{N} \sum_{j=1 \atop t_k \in \{0,1\}} \prod_{j=1}^{N} p_j^{i_j} (1-p_j)^{1-i_j}} \prod_{\substack{s=t_{j_1+1}+t_{j_2+...+1}} \prod_{j=1}^{N} E(p_j^{i_j} (1-p_j)^{1-i_j})}\right)
$$
(15)

since $(\gamma_i, i = 1,2,..,N)$ are independent. Using the following relation.

$$
E(p_j^{i_j}(1-p_j)^{1-i_j}) = \begin{cases} E(p_j) & \text{if } i_j = 1\\ E(1-p_j) & \text{if } i_j = 0 \end{cases}
$$
 (16)

(1 *5)* can be simplified to yield **the** following

$$
P_{F0} = \sum_{s=t_1, t_2, t_3, t_4, t_5, t_6, t_7}^{N} \sum_{i_k \in \{0, 1\}} \left(E(p_j) \right)^{i} \cdot \left(1 - E(p_j) \right)^{i - i} \tag{17}
$$

 $E(p_n)$ for a slow Rayleigh fading channel and noncoherent **FSK** detection can be written **as**

$$
E(p_j) = P_f + (1 - 2 P_f) P_c
$$
 (18)

where P_c is given by (10). (17) shows that the average probability $E(p)$ can be used for the j^{th} link in order to arrive at P_{F0} . If the links are identical, as we have assumed here, then $E(p_n)$ is independent of j, as shown by (18). It can be **seen** that **a similar** result for the detection probability, P_{D0} , is valid.

IV. CONCLUSION

In this paper we considered **a** wireless sensor network, which gathers information **in** order **to make** inference on **a** binary hypothesis. Assuming identical sensors and a counting rule at a fusion center, the exact dependence of the fusion false alarm **and** detection probabilities on the reliability of sensor-to-fusion center link **was** examined. Whereas perfect decision *is* possible in the asymptotic case **of** an infinite set of sensors, for the case of finite number of sensors, depending on the noisyness of the **Iink,** the fusion false alarm probability could increase by several-fold.

ACKNOWLEDGMENT

This **work** was supported in part **by** a grant from Material Technology **Center, SIUC,** Carbondale, **E.**

V. REFERENCES

- [1] D.Li,.D.Wong, H. Hu, and A.M. Sayeed, "Detection, classification, and tracking of targets," IEEE *Signal Process. Magazine,* **March** *2002,* **pp. 17-30.**
- [2] J.F. Chamberland and V. Veeravalli, "Decentralized detection **io** sensor networks," IEEE Trans. *Signal Process.,* Feb. *2003,* pp. **407-416.**
- **[3] R.** Niu, **B.** Chen, and **P.K.** Varshmy, "Decision **fusion** rules in wireless sensor networks **using** fading channel statistics," in 2003 *Conference* on *Infurmation Sciences and System,* The Johns Hopkins University, March **2003.**
- **[4] E,** Chen, R. Jiang, T. Kasetkasem, and **P.K.** Varshaey, "Channel aware decision fusion in wireless sensor networks," IEEE Trans. Signal Process., Dec. 2004, pp. **3454-3458.**
- *[5]* J.F. Chamberland and V. Veeravalli, "Asymptotic results for decentralized detection in power constrained wireless sensor networks," IEEE *Jourul Selecf. Areas Commun.,* **Aug. 2004, pp- 1007-1015.**
- [6] A.R. Reibman and L.W. Nolte, "Optimal fault-tolerant signal detection," IEEE *Trans. ASP,* Jan. 1990, pp. **179-** 180.
- **[7]** A.R. Reibman and L.W. Nolte, "Optimal design and performance *of* distributed signal detection systems with fauIts,"ZEEE *Trans. ASP,* Oct. 1990, pp. **1771-1782.**
- **[SI** S.C.A. Thomopoulos and L. Zhang, "Distributed decision fusion with networking delays and channel errors,"InJom. *Sci.,* Dec. 1992, **pp.** 91-118.
- **[9] C.K,** Sestok, **M.A.** Said, **and** A.V. Oppenheim, "Randomized data selection in detection with applications to distributed signal processing," *Procee.* IEEE, Aug. 2003, **pp. 1184-1 198.**
- [**101** *J.* Proakis, *Digital Communicutions,* McGraw Hill, *NY,* 2002.

Fig 1. Probability of False Alarm/Detection Vs Average Cbannel SNR

Fig 2. Probability of False Alarm/Detection Vs Average Channel SNR

Fig 3. Probability of False Alarm/Detection Vs Average Channel SNR

Fig 4. Probability of False Alarm/Detection Vs Average Channel SNR

306